# Pólya permanent problem: 100 years after 

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Joint work with

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2. Gregor Dolinar, Alexander E. Guterman, Bojan Kuzma, On Gibson barrier for Polya problem, Fundamental and Applied Mathematics, 16(8), 2010, 73-86
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$$
\operatorname{det} A=\sum_{\sigma \in \mathfrak{S}_{n}} \operatorname{sgn}(\sigma) a_{1 \sigma(1)} \cdots a_{n \sigma(n)}
$$

and

$$
\operatorname{per} A=\sum_{\sigma \in \mathfrak{S}_{n}} a_{1 \sigma(1)} \cdots a_{n \sigma(n)}
$$

here $A=\left(a_{i j}\right) \in M_{n}(\mathbb{C}), \mathfrak{S}_{n}$ denotes the set of all permutations of the set $\{1,2, \ldots, n\}$. The value $\operatorname{sgn}(\sigma) \in\{-1,1\}$ is the signum of the permutation $\sigma$.
per is a combinatorial invariant:

$$
\operatorname{per}(P A Q)=\operatorname{per} A
$$

for all permutation matrices $P, Q$

## Some applications of permanent

## Derangements problem

In how many ways can a dance be arranged for $n$ married couples, so that no husband dances with his own wife?

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$$
D_{n}=\operatorname{per}\left(\begin{array}{cccc}
0 & 1 & \cdots & 1 \\
1 & 0 & \cdots & \vdots \\
1 & \cdots & \cdots & 1 \\
1 & \cdots & 1 & 0
\end{array}\right)=\operatorname{per}\left(J_{n}-I_{n}\right)=n!\cdot \sum_{k=0}^{n} \frac{(-1)^{k}}{k!}
$$

## Ménage problem or problème des ménages

In how many ways can $n$ married couples be placed at a round table, so that men and women sit in alternate places and no husband sit on either side of his wife?

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In how many ways can $n$ married couples be placed at a round table, so that men and women sit in alternate places and no husband sit on either side of his wife?

$$
U_{n}=\operatorname{per}\left(\begin{array}{cccccc}
0 & 0 & 1 & \cdots & 1 & 1 \\
1 & 0 & 0 & \cdots & \vdots & \vdots \\
1 & 1 & 0 & \cdots & 1 & 1 \\
1 & 1 & \cdots & \cdots & \cdots & 1 \\
1 & 1 & \dddot{1} & 1 & 0 & 0 \\
0 & 1 & \cdots & 1 & 0
\end{array}\right)=\operatorname{per}\left(J_{n}-I_{n}-P_{n}\right)
$$

$P_{n}$ is a permutation matrix of $(1,2)(2,3) \cdots(n-1, n)(n, 1)$.

## Ménage problem or problème des ménages

In how many ways can $n$ married couples be placed at a round table, so that men and women sit in alternate places and no husband sit on either side of his wife?

Sequence number $A 059375$ in on-line encyclopedia of integer sequences
The first terms:
$12,96,3120,115200,5836320,382072320,31488549120, \ldots$

## Ménage problem or problème des ménages

Formulated in 1891 by Édouard Lucas and independently, a few years earlier, by Peter Guthrie Tait in connection with knot theory

Touchard (1934) derived the formula

$$
U_{n}=2 \cdot n!\sum_{k=0}^{n}(-1)^{k} \frac{2 n}{2 n-k}\binom{2 n-k}{k}(n-k)!
$$

## Latin squares

$S$ is a set, $|S|=n$ usually, $S=\{1,2, \ldots, n\}$
A Latin rectangle on $S$ is an $r \times s$ matrix $A$ with $a_{i j} \in S, a_{i j} \neq a_{i l}$, and $a_{i j} \neq a_{k j}$.
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$n \times n$ Latin rectangle is a Latin square.
Problems: 1. To find the number $L(n, n)$ of Latin squares on $S$
2. To find the number $L(r, n)$ of $r \times n$ Latin rectangles on $S$

## Known facts

1. $L(1, n)=1$
2. $L(2, n)=n!\cdot D_{n}$
3. $L(3, n)=n!\cdot \sum_{k=0}^{\lfloor n / 2\rfloor} C_{n}^{k} D_{n-k} D_{k} U_{n-2 k}$
$\Lambda_{n}^{k}$ is the set of $(0,1)$-matrices with $k 1$ in each row and column.
$m(k, n)$ and $M(k, n)$ are lower and upper bounds for permanent in $\Lambda_{n}^{k}$.
Then

$$
n!D_{n} \prod_{t=2}^{r-1} m(n-t, n) \leq L(r, n) \leq n!D_{n} \prod_{t=2}^{r-1} M(n-t, n)
$$

## Domino tiling

Consider $m \times n$ rectangular chessboard and $2 \times 1$ dominoes.
A tiling is a placement of dominoes that covers all the cells of the board perfectly.


1. If there exists a tiling if we consider a usual chess-board with one corner-cell deleted?
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NO. The total number of cells is odd.

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NO. The total number of cells is odd.
2. If there exists a tiling if we consider a usual chess-board with two opposite corner-cells deleted?

NO. Both deleted cells are of the same color, but domino covers two cells of different colors

## Problems:

1. Existence of tilings.
2. If there are tilings, how many are them?

## Problems:

1. For which $m, n$ do there $\exists$ tilings?
2. If there are tilings, how many are them?

Theorem. Tiling exist $\Leftrightarrow m, n$ are NOT both odd (i.e. $m n$ is even).

Example.


Example.


$$
\begin{gathered}
T(2, n)=T(2, n-1)+T(2, n-2) \\
T(3,2 n)=4 T(3,2 n-2)-T(3,2 n-4)
\end{gathered}
$$

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$$
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\end{gathered}
$$

Difficult recurrent formulas...

Perfect matching in a graph is a selection of edges that covers each vertex exactly once. tilings $\longleftrightarrow$ perfect matchings in underlying grid graph

Chessboard coloring $\Longrightarrow$ bipartite graph
Bipartite graph $\Longrightarrow$ adjacency matrix $A$
The number of tilings $=$ number of perfect matchings $=\operatorname{per}(A)$

The number of tilings: Temperley \& Fisher (1961) and Kasteleyn (1961)

$$
\prod_{j=1}^{m} \prod_{k=1}^{n}\left(4 \cos ^{2} \frac{\pi j}{m+1}+4 \cos ^{2} \frac{\pi k}{n+1}\right)^{\frac{1}{4}}
$$

equivalent to

$$
\prod_{j=1}^{\left\lceil\frac{m}{2}\right\rceil} \prod_{k=1}^{\left\lceil\frac{n}{2}\right\rceil}\left(4 \cos ^{2} \frac{\pi j}{m+1}+4 \cos ^{2} \frac{\pi k}{n+1}\right) .
$$

If $m$ or $n$ is 2: the sequence reduces to the Fibonacci sequence (sequence A000045 in OEIS) (Klarner \& Pollack 1980)

## Applications of permanent:

Counting function for combinatorial problems

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DNA identification

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Makes everybody happy

|  | $\operatorname{det}$ | per |
| :---: | :---: | :---: |
| Geometry | Oriented volume | Combinatorial geometry |
| Algebra | $\lambda_{1} \cdots \lambda_{n}$ | Bounds |
| Complexity | $O\left(n^{3}\right)$ | $\sim(n-1) \cdot\left(2^{n}-1\right)$ |

Ryser's formula

$$
\operatorname{per}(A)=\sum_{t=0}^{n-1}(-1)^{t} \sum_{X \in \Lambda_{n-t}} \prod_{i=1}^{n} r_{i}(X)
$$

$r_{i}(X)=\sum_{j=1}^{t} x_{i j}-i$ th row sum
$\Lambda_{n-t}$ - the set of all $n \times(n-t)$ submatrices of $A$

How many tilings ?
To compute permanent is HARD!
Even if the entries are just 0,1 , computing the permanent is $\sharp P$-complete.

The quantity of transformations preserving a given matrix invariant provide a "measure" of its complexity

Theorem 1 [Frobenius, 1896]

$$
T: M_{n}(\mathbb{C}) \rightarrow M_{n}(\mathbb{C})
$$

- linear, bijective

$$
\operatorname{det}(T(A))=\operatorname{det} A \quad \forall A \in M_{n}(\mathbb{C})
$$

$\Downarrow$

$$
\exists P, Q \in G L_{n}(\mathbb{C}), \operatorname{det}(P Q)=1:
$$

$$
T(A)=P A Q \quad \forall A \in M_{n}(\mathbb{C}) \text { or } T(A)=P A^{t} Q \quad \forall A \in M_{n}(\mathbb{C})
$$

Theorem 2 [Marcus, May] Linear transformation $T$ is permanent pre-
server iff

$$
T(A)=P_{1} D_{1} A D_{2} P_{2} \quad \forall A \in M_{n}(\mathbb{F}), \text { or }
$$

$$
T(A)=P_{1} D_{1} A^{t} D_{2} P_{2} \quad \forall A \in M_{n}(\mathbb{F})
$$

here $D_{i}$ are invertible diagonal matrices, $i=1,2$
$P_{i}$ are permutation matrices, $i=1,2$

Polya, 1913 observed:

$$
n=2:
$$

$$
\operatorname{det}\left(\begin{array}{ll}
a & b \\
c & d
\end{array}\right)=\operatorname{per}\left(\begin{array}{rr}
a & b \\
-c & d
\end{array}\right)
$$

Problem 1. Polya, 1913. Does $\exists$ a uniform way of affixing $\pm$ to the entries of $A=\left(a_{i j}\right) \in M_{n}(\mathbb{F}): \operatorname{per}\left(a_{i j}\right)=\operatorname{det}\left( \pm a_{i j}\right)$ ?

$$
n=2:\left(\begin{array}{ll}
a & b \\
c & d
\end{array}\right) \mapsto\left(\begin{array}{rr}
a & b \\
-c & d
\end{array}\right)
$$

Szegö, 1914. $n>2$ : NO.

## Why NOT ?

$$
n=3: \text { consider } J_{3}=\left(\begin{array}{lll}
1 & 1 & 1 \\
1 & 1 & 1 \\
1 & 1 & 1
\end{array}\right) .
$$

Then per $J_{3}=6$ but

$$
\operatorname{det}\left(\begin{array}{ccc} 
\pm 1 & \pm 1 & \pm 1 \\
\pm 1 & \pm 1 & \pm 1 \\
\pm 1 & \pm 1 & \pm 1
\end{array}\right)<6
$$

since each -1 is in two summands, so all 6 summands can not be positive.

## What about SUBSETS of $M_{n}$ ?

Sometimes the conversion is possible:

1. $\left(\begin{array}{ccc}a & b & 0 \\ c & d & e \\ f & g & h\end{array}\right) \mapsto\left(\begin{array}{ccc}a & b & 0 \\ -c & d & e \\ f-g & h\end{array}\right)$
2. $A: a_{i j}=0$ if $j-i \geq 2$ (Hessenberg matrices)
$A \mapsto \tilde{A}=(\tilde{a i j}): \tilde{a_{i j}}= \begin{cases}-a_{i j}, & \text { if } j-i=1 \\ a_{i j}, & \text { otherwise }\end{cases}$
3. $A$ is Jacobi (3-diagonal) matrix.
$A \mapsto \widehat{A}=\left(\widehat{a_{i j}}\right):$

$$
\widehat{a_{s t}}= \begin{cases}\mathrm{i} a_{s t}, & \text { if } s \neq t \\ a_{s s}, & \text { if } s=\mathrm{t}\end{cases}
$$

Problem 2. Under what conditions does there exist a transformation
$\Phi: M_{n}(\mathbb{F}) \rightarrow M_{m}(\mathbb{F})$ satisfying

$$
\operatorname{per} A=\operatorname{det} \Phi(A) ?
$$

Here a transformation $\Phi$ on $M_{n}(\mathbb{F})$ is called a converter.

Are there linear transformations of this type?

## Are there linear transformations of this type?

Theorem (Marcus, Minc, 1961). There is no bijective linear transformation $\Phi: M_{n}(\mathbb{F}) \rightarrow M_{n}(\mathbb{F}), n>2$ satisfying $\operatorname{per} A=\operatorname{det} \Phi(A) \forall$ $A \in M_{n}(\mathbb{F})$.

Proof: based on linear algebra.

## Are there linear transformations of this type ?

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Proof: based on linear algebra.

Theorem (J. von zur Gathen, 1987). Let $\mathbb{F}$ be infinite, $\operatorname{char}(\mathbb{F}) \neq 2$.
There is no bijective affine transformation $\Phi: M_{n}(\mathbb{F}) \rightarrow M_{n}(\mathbb{F}), n>2$
satisfying per $A=\operatorname{det} \Phi(A) \forall A \in M_{n}(\mathbb{F})$.
Proof: based on algebraic geomery.
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Example. There are non-bijective non-linear converters $\Phi: M_{n}(\mathbb{F}) \rightarrow$
$M_{m}(\mathbb{F})$ of per and det:

$$
\Phi: A \mapsto\binom{1 \frac{1}{2}(\operatorname{det} A-\operatorname{per} A)}{1 \frac{1}{2}(\operatorname{det} A+\operatorname{per} A)} \oplus \operatorname{Id}_{m-2}
$$

Hence, per $A=\operatorname{det} \Phi(A)$ and $\operatorname{det} A=\operatorname{per} \Phi(A)$.

Example. There are bijective non-linear converters of per and det over infinite fields:

For any $\mathbb{F}$ and any $\lambda, \mu \in \mathbb{F}$

$$
\begin{aligned}
& \operatorname{card}\left\{A \in M_{n}(\mathbb{F}) \mid \operatorname{det} A=\mu, \operatorname{per} A=\lambda\right\}= \\
& =\operatorname{card} \mathbb{F} \\
& =\operatorname{card}\left\{A \in M_{n}(\mathbb{F}) \mid \operatorname{det} A=\lambda, \operatorname{per} A=\mu\right\},
\end{aligned}
$$

thus partial bijections exist, and hence the bijection exists.

## WHAT HAPPENS OVER FINITE FIELDS ?

Theorem. [Dolinar, Guterman, Kuzma, Orel] For any $n \geq 3$ there exists $q_{0}=q_{0}(n)$ such that for any finite field $\mathbb{F}$, ch $\mathbb{F} \neq 2,|\mathbb{F}| \geq$ $q_{0}$ there are NO bijective maps $\Phi: M_{n}(\mathbb{F}) \rightarrow M_{n}(\mathbb{F})$ satisfying

$$
\begin{equation*}
\operatorname{per} A=\operatorname{det} \Phi(A) \tag{1}
\end{equation*}
$$

If $n=3$ the conclusion holds for any finite field with ch $\mathbb{F} \neq 2$.

$$
\left|D_{n}\right|=\left|M_{n}\right|-\left|G L_{n}\right|
$$

$\Downarrow$
if $n \geq 4$

$$
\left|D_{n}\right|=q^{n^{2}}-\prod_{k=1}^{n}\left(q^{n}-q^{k-1}\right)=q^{n^{2}-1}+q^{n^{2}-2}+O\left(q^{n^{2}-5}\right)
$$

$$
\left|D_{n}\right|=q^{n^{2}}-\prod_{k=1}^{n}\left(q^{n}-q^{k-1}\right)=q^{n^{2}-1}+q^{n^{2}-2}+0+0+O\left(q^{n^{2}-5}\right)
$$

$$
\begin{aligned}
& L_{n}=q^{n^{2}-1}-q^{n^{2}-2}+O\left(q^{n^{2}-3}\right) \quad(n \geq 4) \\
& U_{n}=q^{n^{2}-1}+0+O\left(q^{n^{2}-3}\right) \quad(n \geq 4)
\end{aligned}
$$

$$
L_{n} \leq P_{n} \leq U_{n}<D_{n}
$$

$\left|P_{n}\right| \leq U_{n}<\left|D_{n}\right|$ if $q$ is sufficiently large ( $q \geq q_{0}$ ).

$$
U_{n}=q^{n^{2}-1}+O\left(q^{n^{2}-3}\right)
$$

$$
\left|D_{n}\right|=q^{n^{2}-1}+q^{n^{2}-2}+O\left(q^{n^{2}-5}\right)
$$

| $n$ | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $q_{0}$ | 3 | 43 | 79 | 121 | 167 | 223 | 289 | 367 | 449 |
| $n$ | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 20 |
| $q_{0}$ | 541 | 641 | 751 | 877 | 997 | 1151 | 1279 | 1433 | 1597 |

## Probability

[P. Erdös, A. Rényi] What is the probability of the permanent of a given matrix to be equal to 0 ?

Theorem. Let $\mathbb{F}$ be a finite field, ch $\mathbb{F} \neq 2 . \forall \lambda \in \mathbb{F}$

$$
P(\text { per } A=\lambda)=\frac{1}{q}+O\left(\frac{1}{q^{2}}\right) .
$$

Let us consider tensor of permanent of $A \in M_{k, n}, k \leq n$ which is
defined by
$T_{A}^{i_{1}, \ldots, i_{n-k}}=\left\{\begin{array}{l}\operatorname{per}\left(A\left(\mid i_{1}, \ldots, i_{n-k}\right)\right), \text { if all } i_{1}, \ldots, i_{n-k} \text { are different } \\ 0, \text { otherwise } .\end{array}\right.$

Examples:

1. $k=n$. Then $T_{A}=\operatorname{per} A$.
2. $k=1, A=\left(a_{1}, \ldots, a_{n}\right)$. Then $T_{A}^{1, \ldots, i-1, i+1, \ldots, n}=a_{i}$.

## Properties:

1. $A \in M_{1, n}$ is a vector. Then $T_{A} \equiv 0$ if and only if $A \equiv 0$.
2. For any $A$ it holds $T_{A}$ is symmetric.

Definition. The convolution of $T_{B}, B \in M_{k, n}$ and $x \in \mathbb{F}_{q}^{n}$ is

$$
\left(T_{B} \circ x\right)^{i_{1}, \ldots, i_{n-k-1}}=\sum_{j=1}^{n} T^{i_{1}, \ldots, i_{n-k-1}, j} \cdot x_{j} \text { of the valency }(n-k-1) .
$$

Lemma. Let $a \in \mathbb{F}_{q}^{n}, A \in M_{k, n}, k<n, B=\binom{a}{A}$. Then $T_{B}=T_{A} \circ a$.

Corollary. For $A \in M_{n}\left(\mathbb{F}_{q}\right)$ formed by the rows $a_{1}, \ldots, a_{n}$.

$$
\begin{gathered}
\operatorname{per}(A)=T_{A}=T_{\left(\begin{array}{c}
a_{1} \\
a_{2} \\
a_{n}
\end{array}\right)=T_{( }^{a_{2}}\left(\begin{array}{c}
a_{3} \\
a_{3} \\
a_{n}
\end{array}\right) \circ a_{1}=\left(T_{\left.\binom{a_{3}}{a_{n}} \circ a_{2}\right) \circ a_{1}=}\right.}^{=\ldots=\left(\ldots\left(T_{a_{n}} \circ a_{n-1}\right) \circ a_{n-2} \ldots\right) \circ a_{1}}
\end{gathered}
$$

Lemma. Let $A \in M_{k \times n}\left(\mathbb{F}_{q}\right)$ and $T_{A} \not \equiv 0$. Then there are at least $q^{n}-q^{k}$ different vectors $x \in F_{q}^{n}$ such that $R=T_{A} \circ x \not \equiv 0$.

Lemma. Let $a=(1, \ldots, 1) \in \mathbb{F}_{q}^{n}, n \geq 3$. Then the number of vectors $x \in \mathbb{F}_{q}^{n}$ such that $R=T_{a} \circ x \neq 0$ is equal to $q^{n}-1>q^{n}-q$.

Theorem (Budrevich, Guterman). Let $\mathbb{F}$ be a finite field, ch $\mathbb{F} \neq 2$.
$\forall n \geq 3$ the number of zeros of per is less
than the number of zeros of det.

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Theorem (Budrevich, Guterman). Let $\mathbb{F}$ be a finite field, ch $\mathbb{F} \neq 2$.
$\forall n \geq 3$ there is $N O$ bijective map $T: M_{n}(\mathbb{F}) \rightarrow M_{n}(\mathbb{F})$ satisfying

$$
\operatorname{per} A=\operatorname{det} T(A)
$$

Problem 3 (Polya). Given a ( 0,1 )-matrix $A \in M_{n}(\mathbb{F})$, does $\exists B$, obtained by changing some of the +1 entries of $A$ into -1 , so that

$$
\text { per } A=\operatorname{det} B ?
$$

The following problems are equivalent to the problem above:

1. Even cycle: A digraph. Does it have no directed circuits of even length?
2. Sign solvability: When does a real square matrix have the property that every real matrix with the same sign pattern is non-singular?

There are more than 30 equivalent problems of this kind, see [ W. Mc-
Cuaig, The Electronic Journal of Combinatorics 11 (2004), R79].

Let $M_{n}$ be the set of all $n \times n\{0,1\}$ matrices over $\mathcal{R}$ - any ring of characteristic 0 .
$\mathrm{S}_{n} \subseteq M_{n}$ - subset of symmetric matrices.
$v(A)$ is the number of 1 of $A$. NB: $v(A)=\sum$ all entries of $A$.
$X \circ A$ denote the Schur (entrywise) product of two matrices.

## Definition.

$A \in M_{n}$ is convertible if $\exists X \in M_{n}( \pm 1)$ :

$$
\operatorname{per} A=\operatorname{det}(X \circ A)
$$

$A \in \mathrm{~S}_{n}$ is symmetrically convertible, if $\exists X \in \mathrm{~S}_{n}( \pm 1)$ :

$$
\operatorname{per} A=\operatorname{det}(X \circ A)
$$

$A \in \mathrm{~S}_{n}$ is symmetrically weakly-convertible, if $\exists X \in \mathrm{~S}_{n}( \pm 1)$ :

$$
\operatorname{per} A=|\operatorname{det}(X \circ A)|
$$

## OBSERVATION

$A \in \mathrm{~S}_{n}$ is symmetrically weakly-convertible.
Then $A$ is convertible.

Multiply a row of $A$ with -1 .

## Can a matrix with arbitrary number of units be convertible ?

Theorem. [Gibson, 1971] Let $A \in M_{n}$ be a convertible matrix with
$\operatorname{per} A>0$.
Then $v(A) \leq \Omega_{n}:=\frac{n^{2}+3 n-2}{2}$.
The equality holds $\Leftrightarrow \exists$ permutation matrices $P, Q: A=P T_{n} Q$.

Here Gibson matrix $T_{n}=\left(t_{i j}\right) \in M_{n}$ is $t_{i j}=\left\{\begin{array}{ll}0, & \text { if } 1 \leq i<j<n \\ 1, & \text { if } i \geq j \text { or } j=n\end{array}\right.$.
Also $G_{n}=\left(g_{i j}\right): \quad g_{i j}= \begin{cases}0, & \text { if } i+j \leq n-1 \\ 1, & \text { if } i+j>n-1\end{cases}$
Note that $G_{n}=T_{n} Q_{n}$ for $Q_{n}=Q(\sigma)$ s.t.

$$
\sigma=(1, n-1)(2, n-2), \ldots,(\lfloor n / 2\rfloor,\lfloor(n+1) / 2\rfloor),
$$

here $\lfloor x\rfloor$ is the largest integer $\leq x$.

$$
T_{5}=\left(\begin{array}{ccccc}
1 & 0 & 0 & 1 \\
1 & 1 & 0 & 0 & 1 \\
1 & 1 & 0 & 1 \\
1 & 1 & 1 & 1 & 1 \\
1 & 1 & 1 & 1 & 1
\end{array}\right) \quad G_{5}=\left(\begin{array}{llllll}
0 & 0 & 0 & 1 & 1 \\
0 & 0 & 1 & 1 & 1 \\
0 & 1 & 1 & 1 & 1 \\
1 & 1 & 1 & 1 & 1 \\
1 & 1 & 1 & 1 & 1
\end{array}\right)
$$

The following results are from [Dolinar, Guterman, Kuzma]

Theorem. $n \geq 3, A \in S_{n}$, per $A>0, A$ is convertible.
Then $v(A) \leq \Omega_{n}=\frac{n^{2}+3 n-2}{2}$.
Let $v(A)=\Omega_{n}$ then $A$ is convertible $\Leftrightarrow A=P G_{n} P^{t}$ for some permutation matrix $P$.

Theorem (Dolinar, Guterman, Kuzma). $n \geq 3, A \in S_{n}$, per $A>0$, $v(A)=\Omega_{n}=\frac{n^{2}+3 n-2}{2}$ and $A$ is convertible. Then
$n \neq 2(\bmod 4) \Longrightarrow A$ is symmetrically convertible.
$n=2(\bmod 4) \Longrightarrow A$ is symmetrically weakly-convertible, but not symmetrically convertible.

Can we find $\omega_{n}$ s.t. $\forall A: v(A)<\omega_{n} \Rightarrow A$ is convertible ?
$\omega_{n}=n+5$

Theorem. [Little, 1972, Graph Theory approach] $n \geq 2, A \in M_{n}$, $v(A) \leq n+5 \Rightarrow A$ is convertible.

## Is $n+5$ a really magic number ?

Theorem (Dolinar, Guterman, Kuzma). $n \geq 3, A \in M_{n}, v(A)=n+6$.
Then $A$ is not convertible $\Leftrightarrow \exists$ permutation matrices $P, Q: P A Q=$ $\operatorname{Id}_{n-3} \oplus J_{3}$, where $J_{3}=\left(\begin{array}{ccc}1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1\end{array}\right)$.

## What is about $S_{n}$ ?

Theorem. [Dolinar, Guterman, Kuzma]

1. $n \geq 2, A \in \mathrm{~S}_{n}, v(A) \leq n+5 \Rightarrow A$ is symmetrically weaklyconvertible.
2. $n \geq 3, v(A)=n+6$. Then $A$ is not convertible $\Leftrightarrow \exists$ permutation matrices $P, Q: P A Q=\operatorname{Id}_{n-3} \oplus J_{3}$.
3. $A$ is convertible, $v(A)=n+6, \Rightarrow A$ is symmetrically weaklyconvertible.

## What happens in between $\omega_{n}$ and $\Omega_{n}$ ?

Theorem. [Dolinar, Guterman, Kuzma] Let $r \in \mathbb{Z}: \omega_{n} \leq r \leq \Omega_{n}$.
Then

1. $\exists A \in \mathrm{~S}_{n}$ : symmetrically weakly-convertible, $\operatorname{per}(A) \neq 0, v(A)=r$
2. $\exists B \in S_{n}$ : not convertible, $v(B)=r$

For fully indecomposable matrices lower bound can be improved.

For fully indecomposable matrices lower bound can be improved.
$A \in M_{n}$ is decomposable, if $\exists$ a permutation matrix $P \in M_{n}$ such that
$A=P\left(\begin{array}{ll}B & 0 \\ C & D\end{array}\right) P^{t}$, where $B, D$ are square.
If $A$ is not decomposable, it is called indecomposable.
$A \in M_{n}$ is partially decomposable if $\exists$ permutation matrices $P, Q \in M_{n}$ such that

$$
A=P\left(\begin{array}{ll}
B & 0 \\
C & D
\end{array}\right) Q
$$

where $B, D$ are square.
If $A$ is no partially decomposable, it is fully indecomposable.
Note, $O$ is decomposable and partially decomposable.

## Lemma.

- If $A \in M_{n}$ is decomposable, then $A$ is partially decomposable.
- If $A \in M_{n}$ is fully indecomposable, then $A$ is indecomposable.

Example.

$$
A=\left(\begin{array}{ll}
0 & 1 \\
1 & 1
\end{array}\right) \in M_{2}
$$

is indecomposable, but partially decomposable with $P=\left(\begin{array}{ll}0 & 1 \\ 1 & 0\end{array}\right), Q=I$.

For fully indecomposable matrices lower bound can be improved.

Theorem (Budrevich, Dolinar, Guterman, Kuzma). Let $A \in M_{n}$ be
fully indecomposable, $v(A) \leq 2 n+2$. Then $A$ is convertible.

Example. Let $A=\left(\begin{array}{lll}1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1\end{array}\right)$. Then $A$ is fully indecomposable, nonconvertible, and $v(A)=9=2 \cdot 3+3$.

Example. Let

$$
A=\left(\begin{array}{ccccccc}
0 & 1 & 1 & 1 & 0 & 0 & \cdots
\end{array}\right)\left(\begin{array}{ccccc}
0 & 1 & 1 & 1 & 0 \\
0 & 1 & 1 & 1 & 0 \\
0 & 0 & \cdots & 0 \\
0 & 0 & 0 & 0 & 1 \\
0 & 0 & \cdots & 0 \\
0 & 0 & \cdots & \cdots & \cdots \\
0 & 0 & \vdots \\
0 & 0 & \cdots & \cdots & \cdots
\end{array}\right)
$$

Then $A$ is not fully indecomposable, indecomposable, not convertible, and $v(A)=n+6$.

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Important trivial observation:

Zeros are better than ones since they are stable under the sign operation!

